Cryptography and Internet Security

How mathematics makes it safe to shop on-line

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http://www.math.unl.edu/~jorr/presentations

Goals

- Bad guys on the net: Why we need internet security
- Codes and ciphers: Julius Caesar and MI5
- The chicken and the egg: Asymmetric ciphers
- 525,600 minutes : Why asymmetric ciphers work
- The bad guys get smart: Man-in-the-middle attacks
- Digital signatures and certificate authorities
- Security ain't safety: Phishing

Bad Guys on the Net Why we need internet security













C:\WINNT\system32\cmd.exe												
Microsoft Windows XP [Version 5.1.2600] (C) Copyright 1985-2001 Microsoft Corp.												
H:\>pathping www.math.unl.edu												
Tracing route to mobius.unl.edu [129.93.180.31] over a maximum of 30 hops: 0 fyb043000009.lancs.local [148.88.169.40] 1 148.88.168.1 2 cp10k-fy5i.rtr.lancs.ac.uk [148.88.255.93] 3 bar?i-cp10k.rtr.lancs.ac.uk [148.88.255.18] 4 194.81.46.1 5 so-1-3-0.warr-sbr1.ja.net [146.97.42.177] 6 so-0-2-0.read-sbr1.ja.net [146.97.33.109] 7 lond-scr3.ja.net [146.97.33.142] 8 po1-0.gn2-gw1.ja.net [146.97.35.98] 9 janet.rt2.lon.uk.geant2.net [62.40.124.197] 10 so-2-0-0.rt1.ams.nl.geant2.net [62.40.112.137] 11 so-7-0-0.rt1.nyc.us.geant2.net [62.40.112.134] 12 198.32.11.50 13 so-0-0-0.0.rtr.wash.net.internet2.edu [64.57.28.11] 14 64.57.28.12 15 64.57.28.0 16 iplsng-chinng.abilene.ucaid.edu [198.32.8.77] 17 kscyng-iplsng.abilene.ucaid.edu [198.32.8.81] 18 ks-2-p00.r.greatplains.net [164.113.238.194] 19 ks-4-t2.r.greatplains.net [164.113.238.206] 20 wsec6-fa-3-45.unl.edu [129.93.5.45]												
Computing statistics for 500 seconds H:\>_												



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 🚰 math.unl.edu - PuTTY
math> telnet www.math.unl.edu 80
Trying 129.93.180.31...
 Connected to www.math.unl.edu.
Escape character is '^]'.
POST /~jorr1/presentations/2007/ea/toaster/form.php HTTP/1.0
Content-Type: application/x-www-form-urlencoded
Content-Length: 62
name=Harry Potter&card=7890 5678 9877 1111&type=Gringotts Bank
HTTP/1.1 200 OK
Date: Mon, 26 Feb 2007 13:39:36 GMT
Server: Apache/2.0.52 (CentOS)
X-Powered-By: PHP/4.3.9
Content-Length: 379
Connection: close
Content-Type: text/html
 <html>
   <head>
     <title>Order Confirmation</title>
     <style>
       @import URL("style.css");
     </style>
  </head>
   <body>
     <h1>Order Confirmation</h1>
    Hi Harry Potter. <br>
     Your card number is 7890 5678 9877 1111. <br>
     It's a Gringotts Bank card.
     <p>
     Your USB Toaster will be winging it's way to you soon...
   </body>
 </html>
Connection closed by foreign host.
math>
```

Codes and Ciphers Julius Caesar and MI5

... if he had anything confidential to say, he wrote it in cipher, that is, by so changing the order of the letters of the alphabet, that not a word could be made out. If anyone wishes to decipher these, and get at their meaning, he must substitute the fourth letter of the alphabet, namely D, for A, and so with the others.



...si qua occultius perferenda erant, per notas scripsit, id est sic structo litterarum ordine, ut nullum uerbum effici posset: quae si qui inuestigare et persequi uelit, quartam elementorum litteram, id est D pro A et perinde reliquas commutet.

Suetonius Life of Julius Caesar, 56

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Real math.unl.edu - PuTTY
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    \langle \mathbf{p} \rangle
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</html>
Connection closed by foreign host.
math>
```

"... substitute the fourth letter of the alphabet, namely D, for A, and so with the others..."

Harry Potter Eduu 🦻 Srwwhu

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 0 1 2

Α	Н											R									Y				
Α	В	С	D	Ε	F	G	Η		J	Κ	L	Μ	Ν	0	Ρ	Q	R	S	Т	U	V	W	X	Y	Ζ
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	0	1	2
D							Κ										U							В	

K D U U B

1234567891011121321344156671879180.1112123456...

Modular Arithmetic

$$a \equiv b \pmod{m}$$

a-b is a multiple of m $10 + 3 \equiv 1 \pmod{12}$ $(10+3)-1=13-1=12=1\times 12$ $2 - 6 \equiv 8 \pmod{12}$ $(2-6)-8 = -4-8 = -12 = -1 \times 12$







Polyalphabetic Cipher

A cipher is a set of rules for encrypting data.

$$E: a \square a+3 \pmod{26}$$

A cipher is **symmetric** if knowledge of the information needed to encrypt also gives you knowledge of how to decrypt.



The Chicken and the Egg Symmetric and asymmetric ciphers





An asymmetric cipher has two parts:

A **public key** k_{public} encrypts A **private key** k_{private} decrypts

Keep the private key secret – give the public key to anyone.





525,600 Minutes

Why asymmetric ciphers work

So:

An **asymmetric cipher**, or **public key cipher**, is one where knowing the information needed to encrypt doesn't help you decrypt.

How is this possible?

In fact, k_{public} and $k_{\text{private}} \underline{are}$ related, but...



RSA Public Key Cryptography

Described by Rob Rivest, Adi Shamir, and Leonard Adleman at MIT in 1977.

The idea is based on prime numbers...

A prime number is one whose only factors are 1 and itself.

e.g. 2, 3, 5, 7, 11, 13 but not 4, or 6

Theorem. *Every number is the product of prime numbers.* e.g. 1,386 = 2×693 = 2×3×231 = 2×3×3×77=2×3×3×7×11

Theorem. There is no biggest prime number.

If 2,3,5,7,...,*P* were all the prime numbers then what about $1 + 2 \times 3 \times 5 \times 7 \times ... \times P$

Each of these numbers is the product of *exactly* **two** prime numbers. What are they?

$$6 = 2 \times 3$$

$$10 = 2 \times 5$$

$$21 = 3 \times 7$$

$$221 = 13 \times 17$$

$$713 = 23 \times 31$$

$$456,989,977,669 = 611,953 \times 746,773$$

$$= P_{5000} \times P_{6000}$$

The RSA **public** key consists of a number which is the product of two prime numbers. If you could figure out *which* two prime numbers you could find the **private** key.



"Ask a computer – computers are good at these kind of things..."

Look again at

456,989,977,669= 611,953 × 746,773

One way to factor 456,989,977,669 is to check all the numbers 1,2,3,... up to $\sqrt{456,989,977,669} \approx 676,010$.

If a computer can do 1,000,000 tests in a second, then it can do this in just $676,010 \div 1,000,000 = 0.676$ seconds.

But what if N = P×Q is 100 digits long?

Then

$$10^{99} \le N < 10^{100}$$

SO

$$10^{49} \le \sqrt{N} < 10^{50}$$

and the computer can solve it in

$$10^{50} \div 10^{6} = 10^{44}$$
 seconds.



Age of the universe = 13,700,000,000 years



Theorem. There is no biggest prime number.

<u>And</u> we have good algorithms for finding very big prime numbers (100's of digits)

<u>But</u> we have no methods of finding the prime factors of N=PQ that are *qualitatively* better than just checking all possibilities:

$$T = C A^d$$
 where $d = #$ digits in N

How does RSA work?

Need to generate a **public** and a **private** key.

Step 1: Pick two (very) big prime numbers p and q

Step 2: Pick a number 0 < r < (p - 1)(q - 1)

Step 3: Find a number 0 < s < (p - 1)(q - 1) such that

$$rs \equiv 1 \pmod{(p-1)(q-1)}$$

Key Fact: For any pumber of ma - b is a multiple of (p - 1)(q - 1)



$$x^{rs} \equiv x \pmod{pq}$$

$$rs \equiv 1 \pmod{(p-1)(q-1)}$$

$$rs - 1 = k(p-1)(q-1)$$

$$rs = 1 + k(p-1)(q-1)$$

$$x^{rs} = x^{1+k(p-1)(q-1)} = x \cdot (x^{(p-1)(q-1)})^{k}$$

= x \cdot (1)^{k} = x \left(mod pq \right)

 $x^{p-1} \equiv 1 \pmod{p}$ "Fermat's Little Theorem" $x^{q-1} \equiv 1 \pmod{q}$

 $x^{(p-1)(q-1)} \equiv 1 \pmod{pq}$ "Chinese Remainder Theorem"

The Bad Guys Get Smart Man-in-the-middle attacks





Digital Signatures



Anyone can read the message...

...but only **one** person could have <u>written</u> the message...

Bob!





Digital Signatures



Security Ain't Safety Phishing





http://www.math.unl.edu/~jorr/presentations